Have a Pick-up Game

Suppose you are playing a game with another person. The rules are as follows:

- There are *a* objects on the table.
- Both players take turns picking objects from the table. In each turn, a player can take $1, 2, \cdots$ or b objects for some positive integer b.
- The player taking the final object wins.

The goal, naturaly, is to come up with a strategy which guarantees victory. The good news is that there is one.

Let $c = a \pmod{b+1}$. Assuming you pick first, during the first turn, you should pick c objects. For the rest of the game, if the other player picks d objects during a turn, you should pick b+1-d objects in the next turn. (So if the other player picks 2 objects in the last turn, you should pick b-1 objects.)

However, if a is divisible by b + 1 (or c = 0), the player picking second would win by picking b + 1 - d objects immediately after the other player picks d objects. If you pick first you would have to wait until the number of objects remaining is not divisible by b + 1 after the other player's turn.

If the player taking the final object loses and you pick first, you should pick c-1 objects in the first turn. Then apply the above strategy for the rest of the game. (If c = 1, you should pick second.)

From the above, we can see that you should pick second when c = 0 and the player taking the final object wins; or when c = 1 and the player taking the final object loses. Otherwise you should pick first.

Suppose the player picking the final object wins. Since $c = a \pmod{b+1}$, there are a = k(b+1) + c objects at the beginning of the game for some positive integer k. Assume you pick first and you pick $c \neq 0$ objects during the first turn, then there would be k(b+1) objects remaining. During the next turn, the other player would pick d objects and you would pick b+1-d objects. So d+(b+1-d) = b+1 objects would be removed and there would be (k-1)(b+1) objects remaining on the table (after your second turn). By following the above strategy, there would be (k-i)(b+1) objects remaining on the table after your i + 1-th turn. Thus, no objects would remain on the table after your k + 1-th turn, handing you the victory.

If c = 0, you are assured victory by picking second. During the *i*-th turn, the other player picks d objects and you pick b+1-d objects, leaving (k-i)(b+1) objects on the table. So you are able to remove all objects left on the table during your *k*-th term, thus winning the game.

Now suppose the player picking the final object loses and $c \neq 1$. If you pick c-1 objects in the first turn, there would be k(b+1) + 1 objects remaining. During the next turn, the other player would pick d objects and you would pick b+1-d objects. So d+(b+1-d) = b+1 objects would be removed and there would be (k-1)(b+1) + 1 objects remaining on the table (after your second turn). By following the above strategy, there would be (k-i)(b+1)+1 objects remaining on the table after your i + 1-th turn. Thus, 1 object would be left on the table after your k + 1-th turn. The other player would have to pick up the object, giving you the victory.

If c = 1, you are assured victory by picking second. During the *i*-th turn, the other player picks *d* objects and you pick b+1-d objects, leaving (k-i)(b+1)+1 objects on the table. By following this strategy, there would be 1 object left on the table after your *k*-th term, securing the win.